MCM36100

1M x 36 Bit Dynamic Random Access Memory Module

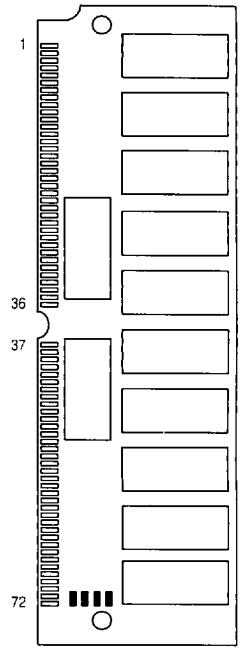
The MCM36100 is a 36M dynamic random access memory (DRAM) module organized as 1,048,576 x 36 bits. The module is a 72-lead single-in-line memory module (SIMM) consisting of eight MCM54400A DRAMs housed in standard 300 mil SOJ packages and four CMOS 1M x 1 DRAMs housed in 20/26 lead SOJ packages, mounted on a substrate along with a 0.22 µF (min) decoupling capacitor mounted under each DRAM. The MCM54400A is a CMOS high-speed dynamic random access memory organized as 1,048,576 four-bit words and fabricated with CMOS silicon-gate process technology.

- Three-State Data Output
- Early-Write Common I/O Capability
- Fast Page Mode Capability
- TTL-Compatible Inputs and Outputs
- $\overline{RAS}$-Only Refresh
- $\overline{CAS}$ Before $\overline{RAS}$ Refresh
- Hidden Refresh
- 1024 Cycle Refresh: 16 ms (Max)
- Consists of Eight 1M x 4 DRAMs, Four 1M x 1 DRAMs, and Twelve 0.22 µF (Min) Decoupling Capacitors
- Unlatched Data Out at Cycle End Allows Two Dimensional Chip Selection
- Fast Access Time (t_{RAC}): MCM36100-60 = 60 ns (Max)
  MCM36100-70 = 70 ns (Max)
- Low Active Power Dissipation: MCM36100-60 = 7.26 W (Max)
  MCM36100-70 = 6.16 W (Max)
- Low Standby Power Dissipation: TTL Levels = 132 mW (Max)
  CMOS Levels = 66 mW (Max)
- Available in Low Profile (1 Inch) Module (MCM36100ASH)

PIN OUT

Pin	Name	Pin	Name	Pin	Name	Pin	Name	Pin	Name	Pin	Name
1	V_{SS}	13	A1	25	DQ24	37	DQ17	49	DQ9	61	DQ14
2	DQ0	14	A2	26	DQ7	38	DQ35	50	DQ27	62	DQ33
3	DQ18	15	A3	27	DQ25	39	V_{SS}	51	DQ10	63	DQ15
4	DQ1	16	A4	28	A7	40	$\overline{CAS0}$	52	DQ28	64	DQ34
5	DQ19	17	A5	29	NC	41	$\overline{CAS2}$	53	DQ11	65	DQ16
6	DQ2	18	A6	30	V_{CC}	42	$\overline{CAS3}$	54	DQ29	66	NC
7	DQ20	19	NC	31	A8	43	$\overline{CAS1}$	55	DQ12	67	PD1
8	DQ3	20	DQ4	32	A9	44	$\overline{RAS0}$	56	DQ30	68	PD2
9	DQ21	21	DQ22	33	NC	45	NC	57	DQ13	69	PD3
10	V_{CC}	22	DQ5	34	$\overline{RAS2}$	46	NC	58	DQ31	70	PD4
11	NC	23	DQ23	35	DQ26	47	$\overline{W}$	59	V_{CC}	71	NC
12	A0	24	DQ6	36	DQ8	48	NC	60	DQ32	72	V_{SS}

S PACKAGE
SIMM MODULE
CASE 866A-02



PIN NAMES

A0 – A9	Address Inputs
DQ0 – DQ35	Data Input/Output
$\overline{CAS0}$ – $\overline{CAS3}$	Column Address Strobe
PD1 – PD4	Presence Detect
$\overline{RAS0}$, $\overline{RAS2}$	Row Address Strobe
$\overline{W}$	Read/Write Input
V_{CC}	Power (+ 5 V)
V_{SS}	Ground
NC	No Connection

All power supply and ground pins must be connected for proper operation of the device.

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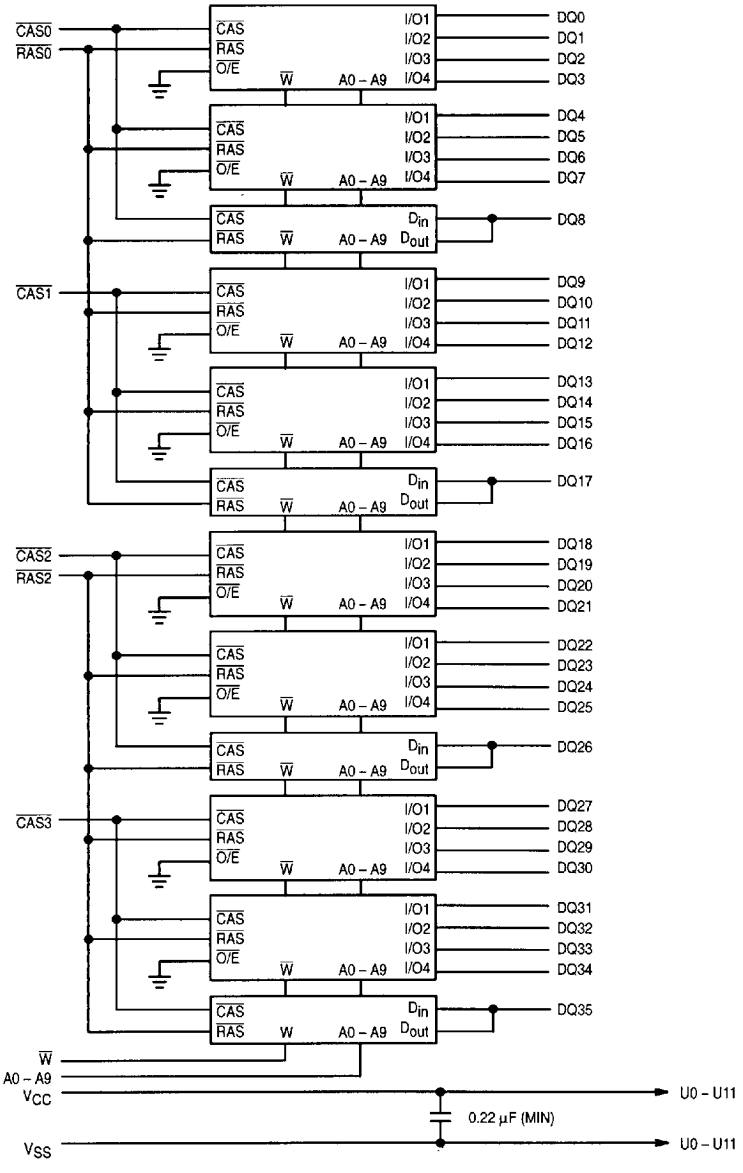
BLOCK DIAGRAM



PRESENCE DETECT PIN OUT		
Pin Name	**60 ns**	**70 ns**
PD1	V_{SS}	V_{SS}
PD2	V_{SS}	V_{SS}
PD3	NC	V_{SS}
PD4	NC	NC

ABSOLUTE MAXIMUM RATINGS (See Note)

Rating	Symbol	Value	Unit
Power Supply Voltage	V_{CC}	– 1 to + 7	V
Voltage Relative to V_{SS} (For Any Pin Except V_{CC})	V_{in}, V_{out}	– 1 to + 7	V
Data Output Current per DQ Pin	I_{out}	50	mA
Power Dissipation	P_D	8.4	W
Operating Temperature Range	T_A	0 to + 70	°C
Storage Temperature Range	T_{stg}	– 55 to + 125	°C

NOTE: Permanent device damage may occur if ABSOLUTE MAXIMUM RATINGS are exceeded. Functional operation should be restricted to RECOMMENDED OPERATING CONDITIONS. Exposure to higher than recommended voltages for extended periods of time could affect device reliability.

This device contains circuitry to protect the inputs against damage due to high static voltages or electric fields; however, it is advised that normal precautions be taken to avoid application of any voltage higher than maximum rated voltages to these high-impedance circuits.

DC OPERATING CONDITIONS AND CHARACTERISTICS

(V_{CC} = 5.0 V ± 10%, T_A = 0 to 70°C, Unless Otherwise Noted)

RECOMMENDED OPERATING CONDITIONS (All voltages referenced to V_{SS})

Parameter	Symbol	Min	Typ	Max	Unit
Supply Voltage (Operating Voltage Range)	V_{CC}	4.5	5.0	5.5	V
	V_{SS}	0	0	0	
Logic High Voltage, All Inputs	V_{IH}	2.4	—	6.5	V
Logic Low Voltage, All Inputs	V_{IL}	– 1.0	—	0.8	V

DC CHARACTERISTICS AND SUPPLY CURRENTS

Characteristic		Symbol	Min	Max	Unit	Notes
V_{CC} Power Supply Current	MCM36100-60, t_{RC} = 110 ns MCM36100-70, t_{RC} = 130 ns	I_{CC1}	— —	1320 1120	mA	1
V_{CC} Power Supply Current (Standby) ($\overline{RAS}$ = $\overline{CAS}$ = V_{IH})		I_{CC2}	—	24	mA	
V_{CC} Power Supply Currentt During $\overline{RAS}$-Only Refresh Cycles	MCM36100-60, t_{RC} = 110 ns MCM36100-70, t_{RC} = 130 ns	I_{CC3}	— —	1320 1120	mA	1
V_{CC} Power Supply Currentt During Fast Page Mode Cycle	MCM36100-60, t_{PC} = 45 ns MCM36100-70, t_{PC} = 45 ns	I_{CC4}	— —	800 800	mA	1, 2
V_{CC} Power Supply Current (Standby) ($\overline{RAS}$ = $\overline{CAS}$ = V_{CC} – 0.2 V)		I_{CC5}	—	12	mA	
V_{CC} Power Supply Current During $\overline{CAS}$ Before $\overline{RAS}$ Refresh Cycle	MCM36100-60, t_{RC} = 110 ns MCM36100-70, t_{RC} = 130 ns	I_{CC6}	— —	1320 1120	mA	1
Input Leakage Current ($V_{SS} \le V_{in} \le V_{CC}$)		$I_{lkg(I)}$	– 120	+ 120	µA	
Output Leakage Current ($\overline{CAS}$ at Logic 1, $V_{SS} \le V_{out} \le V_{CC}$)		$I_{lkg(O)}$	– 10	+ 10	µA	
Output High Voltage (I_{OH} = – 5 mA)		V_{OH}	2.4	—	V	
Output Low Voltage (I_{OL} = 4.2 mA)		V_{OL}	—	0.4	V	

NOTES:
1. Current is a function of cycle rate and output loading; maximum current is measured at the fastest cycle rate with the output open.
2. Measured with one address transition per page mode cycle.

CAPACITANCE (f = 1.0 MHz, T_A = 25°C, V_{CC} = 5 V, Periodically Sampled Rather Than 100% Tested)

Characteristic		Symbol	Min	Max	Unit
Input Capacitance	A0 – A9	C_{I1}	—	70	pF
	$\overline{W}$	C_{I2}		94	
	$\overline{RAS0}$, $\overline{RAS2}$	C_{I3}		52	
	$\overline{CAS0}$ – $\overline{CAS3}$	C_{I4}		31	
I/O Capacitance	DQ0 – DQ7, DQ9 – DQ16, DQ18 – DQ25, DQ27 – DQ34	C_{DQ1}	—	17	pF
	DQ8, DQ17, DQ26, DQ35	C_{DQ2}		22	

NOTE: Capacitance measured with a Boonton Meter or effective capacitance calculated from the equation: C = I Δt/ΔV.

AC OPERATING CONDITIONS AND CHARACTERISTICS

(V_{CC} = 5.0 V ± 10%, T_A = 0 to 70°C, Unless Otherwise Noted)

READ AND WRITE CYCLES (See Notes 1, 2, 3, and 4)

Parameter	Symbol		MCM36100-60		MCM36100-70		Unit	Notes
	Std	Alt	Min	Max	Min	Max		
Random Read or Write Cycle Time	t_{RELREL}	t_{RC}	110	—	130	—	ns	5
Fast Page Mode Cycle Time	t_{CELCEL}	t_{PC}	45	—	45	—	ns	
Access Time from $\overline{RAS}$	t_{RELQV}	t_{RAC}	—	60	—	70	ns	6, 7
Access Time from $\overline{CAS}$	t_{CELQV}	t_{CAC}	—	20	—	20	ns	6, 8
Access Time from Column Address	t_{AVQV}	t_{AA}	—	30	—	35	ns	6, 9
Access Time from Precharge $\overline{CAS}$	t_{CEHQV}	t_{CPA}	—	40	—	40	ns	6
$\overline{CAS}$ to Output in Low-Z	t_{CELQX}	t_{CLZ}	0	—	0	—	ns	6
Output Buffer and Turn-Off Delay	t_{CEHQZ}	t_{OFF}	0	20	0	20	ns	10
Transition Time (Rise and Fall)	t_T	t_T	3	50	3	50	ns	
$\overline{RAS}$ Precharge Time	t_{REHREL}	t_{RP}	40	—	50	—	ns	
$\overline{RAS}$ Pulse Width	t_{RELREH}	t_{RAS}	60	10 k	70	10 k	ns	
$\overline{RAS}$ Pulse Width (Fast Page Mode)	t_{RELREH}	t_{RASP}	60	100 k	70	100 k	ns	
$\overline{RAS}$ Hold Time	t_{CELREH}	t_{RSH}	20	—	20	—	ns	
$\overline{CAS}$ Hold Time	t_{RELCEH}	t_{CSH}	60	—	70	—	ns	
$\overline{CAS}$ Precharge to $\overline{RAS}$ Hold Time	t_{CEHREH}	t_{RHCP}	40	—	40	—	ns	
$\overline{CAS}$ Pulse Width	t_{CELCEH}	t_{CAS}	20	10 k	20	10 k	ns	
$\overline{RAS}$ to $\overline{CAS}$ Delay Time	t_{RELCEL}	t_{RCD}	20	40	20	50	ns	11

NOTES: (continued)

1. V_{IH} (min) and V_{IL} (max) are reference levels for measuring timing of input signals. Transition times are measured between V_{IH} and V_{IL}.
2. An initial pause of 200 µs is required after power-up followed by 8 $\overline{RAS}$ cycles before proper device operation is guaranteed.
3. The transition time specification applies for all input signals. In addition to meeting the transition rate specification, all input signals must transition between V_{IH} and V_{IL} (or between V_{IL} and V_{IH}) in a monotonic manner.
4. AC measurements t_T = 5.0 ns.
5. The specification for t_{RC} (min) is used only to indicate cycle time at which proper operation over the full temperature range (0°C ≤ T_A ≤ 70°C) is ensured.
6. Measured with a current load equivalent to 2 TTL (– 200 µA, + 4 mA) loads and 100 pF with the data output trip points set at V_{OH} = 2.0 V and V_{OL} = 0.8 V.
7. Assumes that $t_{RCD} \leq t_{RCD}$ (max).
8. Assumes that $t_{RCD} \geq t_{RCD}$ (max).
9. Assumes that $t_{RAD} \geq t_{RAD}$ (max).
10. t_{OFF} (max) defines the time at which the output achieves the open circuit condition and is not referenced to output voltage levels.
11. Operation within the t_{RCD} (max) limit ensures that t_{RAC} (max) can be met. t_{RCD} (max) is specified as a reference point only; if t_{RCD} is greater than the specified t_{RCD} (max) limit, then access time is controlled exclusively by t_{CAC}.

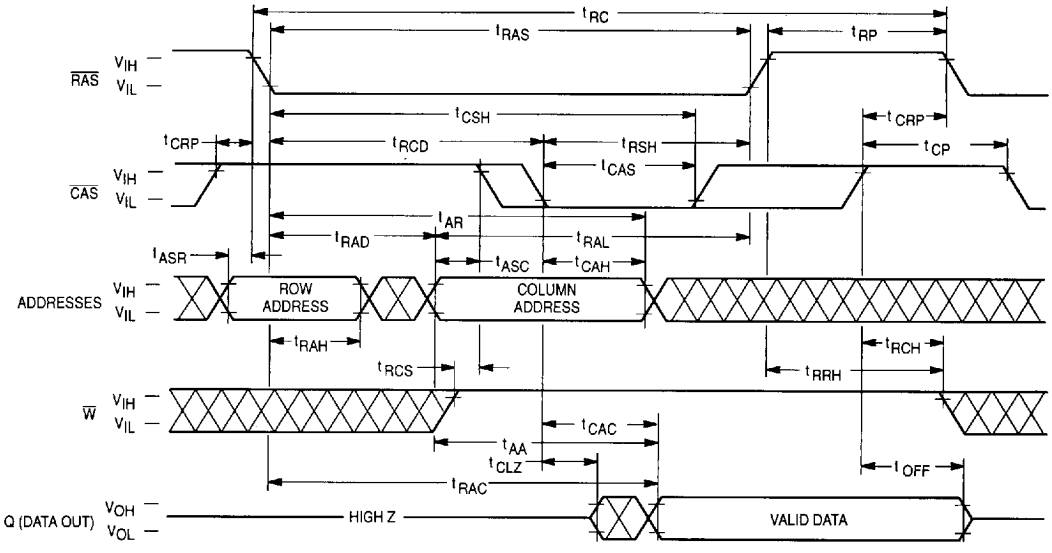
READ AND WRITE CYCLES (Continued)

Parameter	Symbol		MCM36100-60		MCM36100-70		Unit	Notes
	Std	Alt	Min	Max	Min	Max		
$\overline{RAS}$ to Column Address Delay Time	t_{RELAV}	t_{RAD}	15	30	15	35	ns	12
$\overline{CAS}$ to $\overline{RAS}$ Precharge Time	t_{CEHREL}	t_{CRP}	5	—	5	—	ns	
$\overline{CAS}$ Precharge Time	t_{CEHCEL}	t_{CP}	10	—	10	—	ns	
Row Address Setup Time	t_{AVREL}	t_{ASR}	0	—	0	—	ns	
Row Address Hold Time	t_{RELAX}	t_{RAH}	10	—	10	—	ns	
Column Address Setup Time	t_{AVCEL}	t_{ASC}	0	—	0	—	ns	
Column Address Hold Time	t_{CELAX}	t_{CAH}	15	—	15	—	ns	
Column Address to $\overline{RAS}$ Lead Time	t_{AVREH}	t_{RAL}	30	—	35	—	ns	
Read Command Setup Time	t_{WHCEL}	t_{RCS}	0	—	0	—	ns	
Read Command Hold Time Referenced to $\overline{CAS}$	t_{CEHWX}	t_{RCH}	0	—	0	—	ns	13
Read Command Hold Time Referenced to $\overline{RAS}$	t_{REHWX}	t_{RRH}	0	—	0	—	ns	13
Write Command Hold Time Referenced to $\overline{CAS}$	t_{CELWH}	t_{WCH}	10	—	15	—	ns	
Write Command Pulse Width	t_{WLWH}	t_{WP}	10	—	15	—	ns	
Write Command to $\overline{RAS}$ Lead Time	t_{WLREH}	t_{RWL}	20	—	20	—	ns	
Write Command to $\overline{CAS}$ Lead Time	t_{WLCEH}	t_{CWL}	20	—	20	—	ns	
Data in Setup Time	t_{DVCEL}	t_{DS}	0	—	0	—	ns	14
Data in Hold Time	t_{CELDX}	t_{DH}	15	—	15	—	ns	14
Refresh Period	t_{RVRV}	t_{RFSH}	—	16	—	16	ms	
Write Command Setup Time	t_{WLCEL}	t_{WCS}	0	—	0	—	ns	15
$\overline{CAS}$ Setup Time for $\overline{CAS}$ Before $\overline{RAS}$ Refresh	t_{RELCEL}	t_{CSR}	5	—	10	—	ns	
$\overline{CAS}$ Hold Time for $\overline{CAS}$ Before $\overline{RAS}$ Refresh	t_{RELCEH}	t_{CHR}	15	—	30	—	ns	
$\overline{RAS}$ Precharge to $\overline{CAS}$ Active Time	t_{REHCEL}	t_{RPC}	0	—	0	—	ns	
$\overline{CAS}$ Precharge Time for $\overline{CAS}$ Before $\overline{RAS}$ Counter Time	t_{CEHCEL}	t_{CPT}	30	—	40	—	ns	
Write to $\overline{RAS}$ Precharge Time ($\overline{CAS}$ Before $\overline{RAS}$ Refresh)	t_{WHREL}	t_{WRP}	10	—	10	—	ns	
Write to $\overline{RAS}$ Hold Time ($\overline{CAS}$ Before $\overline{RAS}$ Refresh)	t_{RELWL}	t_{WRH}	10	—	10	—	ns	

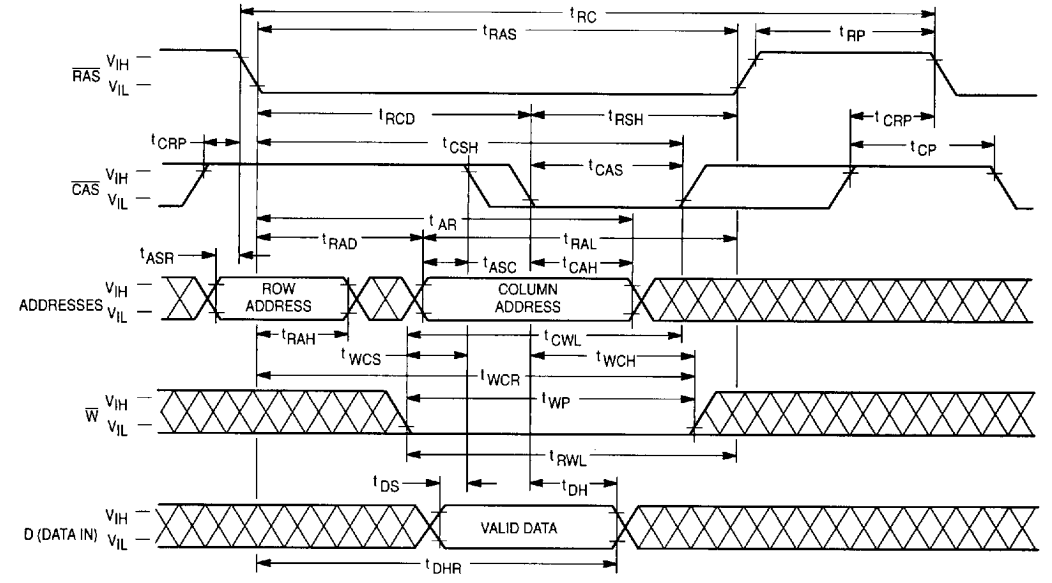
NOTES:

12. Operation within the t_{RAD} (max) limit ensures that t_{RAC} (max) can be met. t_{RAD} (max) is specified as a reference point only; if t_{RAD} is greater than the specified t_{RAD} (max), then access time is controlled exclusively by t_{AA}.
13. Either t_{RRH} or t_{RCH} must be satisfied for a read cycle.
14. These parameters are referenced to $\overline{CAS}$ leading edge in random write cycles.
15. t_{WCS} is not a restrictive operating parameter. It is included in the data sheet as an electrical characteristic only; if $t_{WCS} \geq t_{WCS}$ (min), the cycle is an early write cycle and the data out pin will remain open circuit (high impedance) throughout the entire cycle. If this condition is not satisifed, the condition of the data out (at access time) is indeterminate.

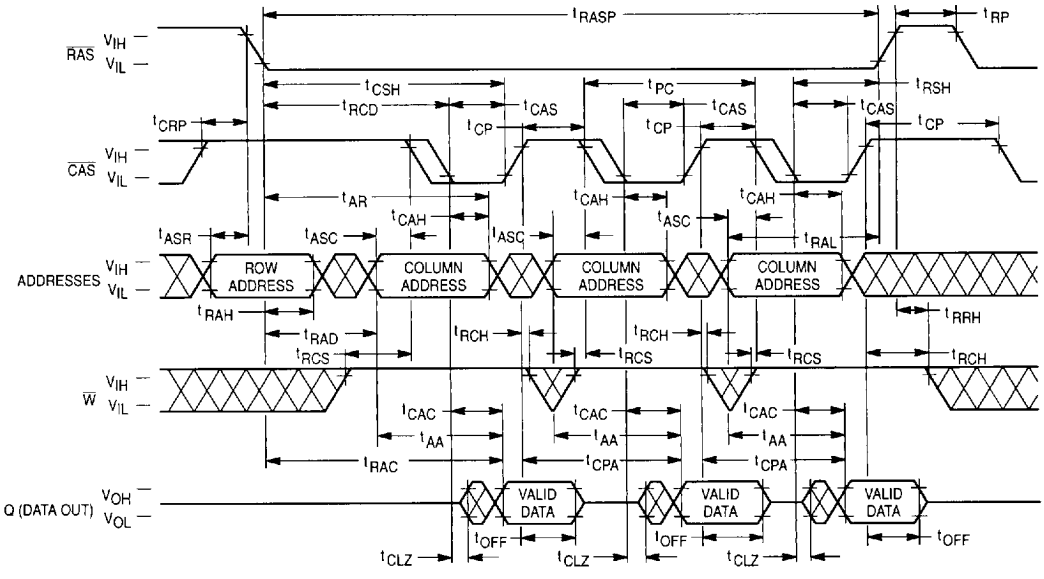
READ CYCLE



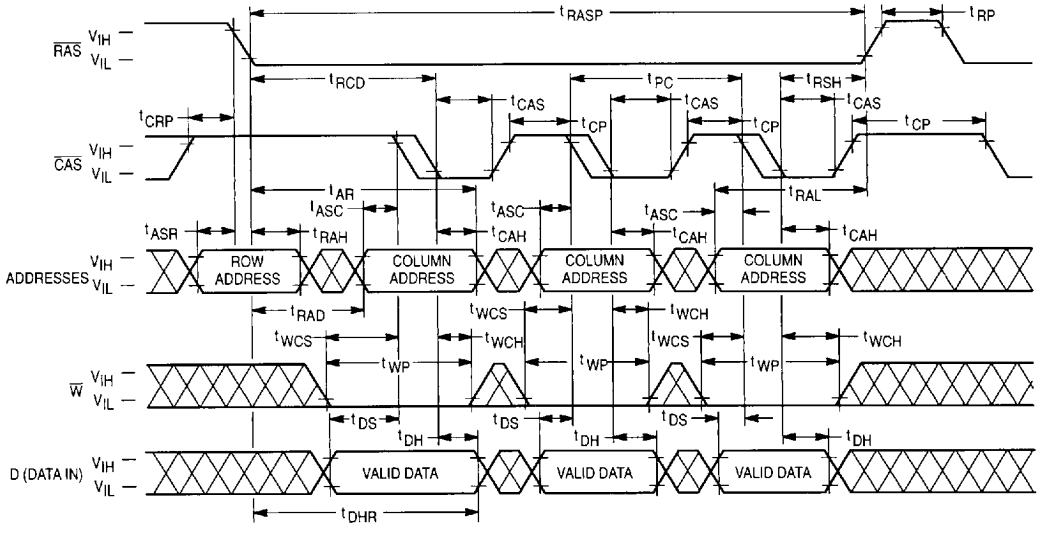
EARLY WRITE CYCLE



FAST PAGE MODE READ CYCLE



FAST PAGE MODE WRITE CYCLE (EARLY WRITE)



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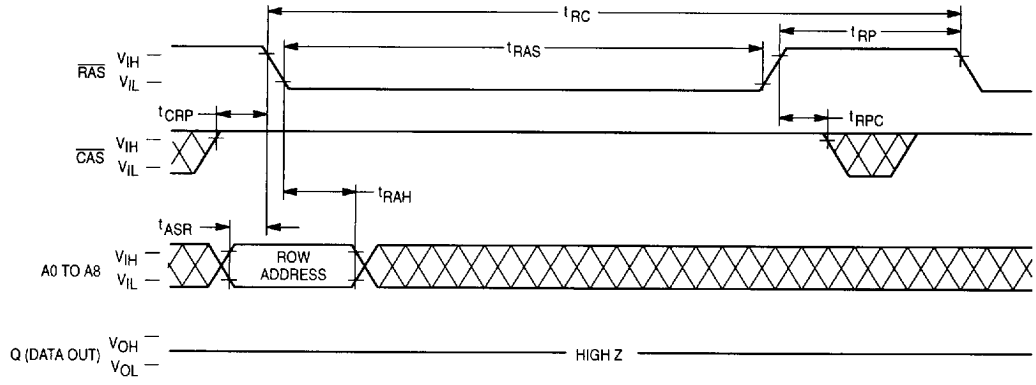
$\overline{RAS}$-ONLY REFRESH CYCLE

($\overline{W}$ and A9 are Don't Care)



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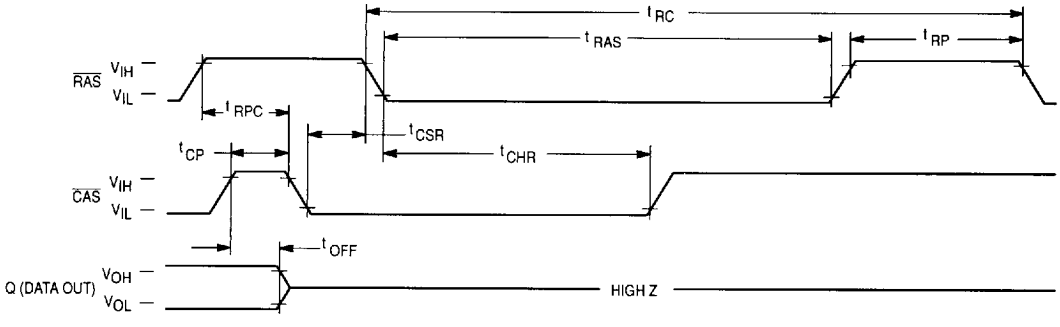
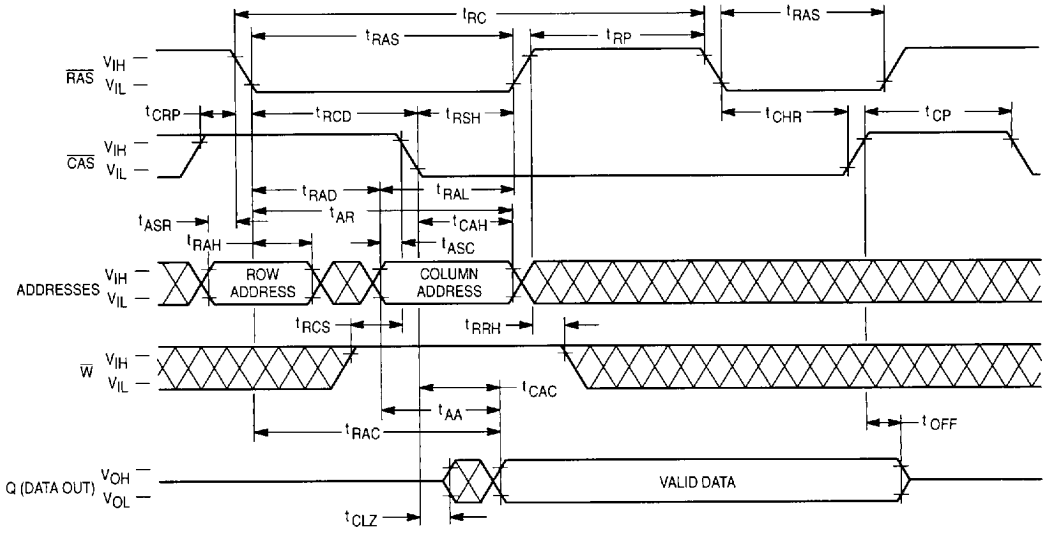
$\overline{CAS}$ BEFORE $\overline{RAS}$ REFRESH CYCLE

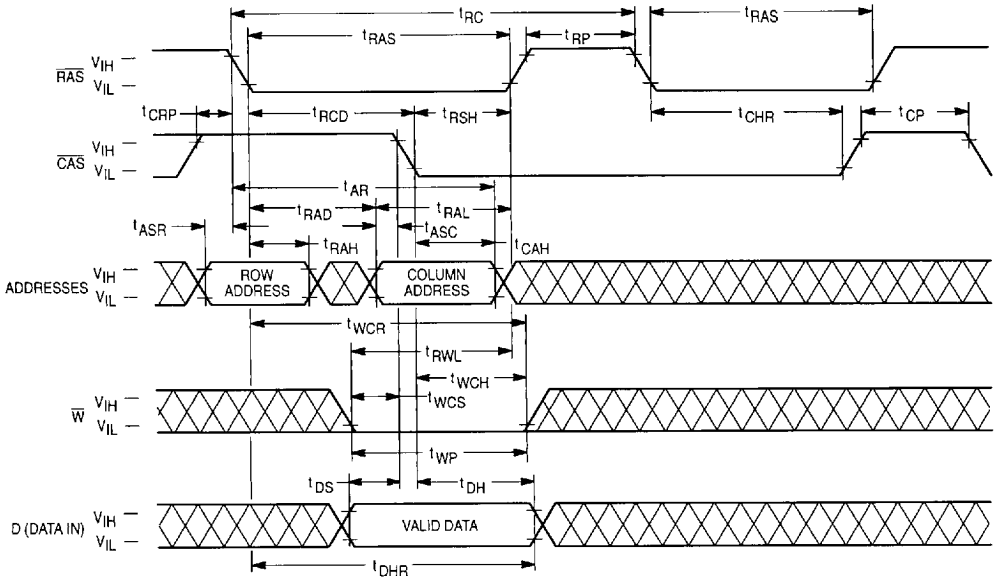
(A0 – A9 are Don't Care)



HIDDEN REFRESH CYCLE (READ)



HIDDEN REFRESH CYCLE (WRITE)



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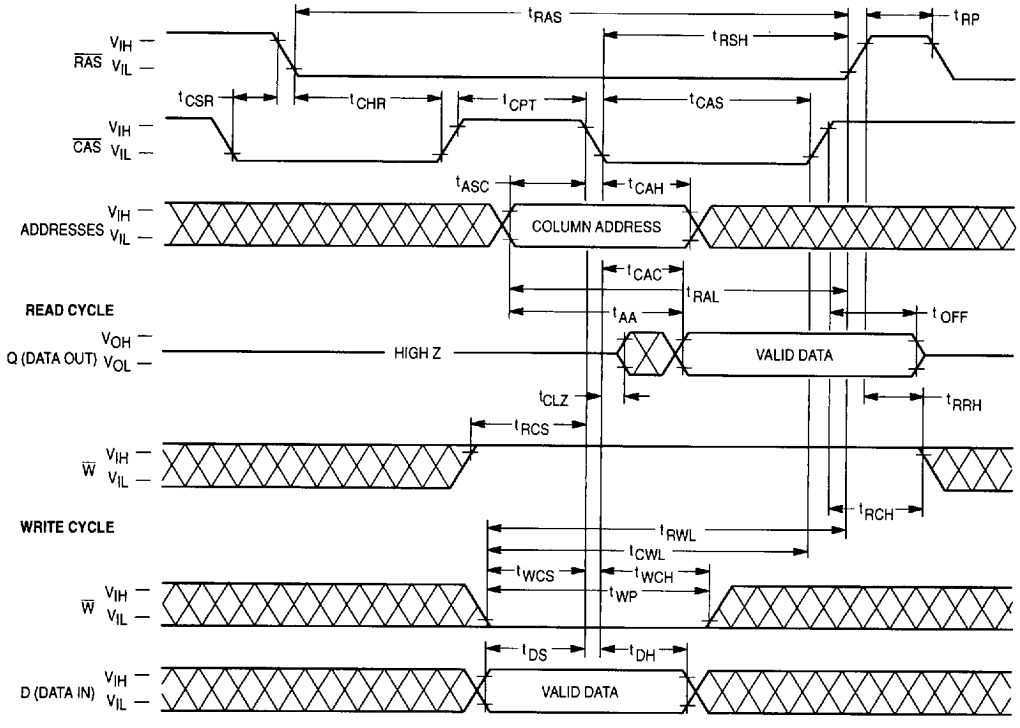
$\overline{CAS}$ BEFORE $\overline{RAS}$ REFRESH COUNTER TEST CYCLE



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DEVICE INITIALIZATION

On power-up, an initial pause of 200 microseconds is required for the internal substrate generator to establish the correct bias voltage. This must be followed by a minimum of eight active cycles of the row address strobe (clock) to initialize all dynamic nodes within the module. During an extended inactive state (greater than 4 milliseconds with the device powered up), a wake up sequence of eight active cycles is necessary to ensure proper operation.

ADDRESSING THE RAM

The ten address pins on the device are time multiplexed at the beginning of a memory cycle by two clocks, row address strobe ($\overline{RAS}$) and column address strobe ($\overline{CAS}$), into two separate 10-bit address fields. A total of twenty address bits, ten rows and ten columns, will decode one of the 1,048,576 word locations in the device. $\overline{RAS}$ active transition is followed by $\overline{CAS}$ active transition (active = V_{IL}, t_{RCD} minimum) for all read or write cycles. The delay between $\overline{RAS}$ and $\overline{CAS}$ active transitions, referred to as the **multiplex window**, gives a system designer flexibility in setting up the external addresses into the RAM.

The external $\overline{CAS}$ signal is ignored until an internal $\overline{RAS}$ signal is available. This "gate" feature on the external $\overline{CAS}$ clock enables the internal $\overline{CAS}$ line as soon as the row address hold time (t_{RAH}) specification is met (and defines t_{RCD} minimum). The multiplex window can be used to absorb skew delays in switching the address bus from row to column addresses and in generating the $\overline{CAS}$ clock.

There are three other variations in addressing the module: **$\overline{RAS}$-only refresh cycle, $\overline{CAS}$ before $\overline{RAS}$ refresh cycle,** and **page mode**. All three are discussed in separate sections that follow.

READ CYCLE

The DRAM may be read with either a "normal" random read cycle or a page mode read cycle. The normal read cycle is outlined here, while the page mode cycle is discussed in a separate section.

The normal read cycle begins as described in **ADDRESSING THE RAM**, with $\overline{RAS}$ and $\overline{CAS}$ active transitions latching the desired bit location. The write ($\overline{W}$) input level must be high (V_{IH}), t_{RCS} (minimum) before the $\overline{CAS}$ active transition, to enable read mode.

Both the $\overline{RAS}$ and $\overline{CAS}$ clocks trigger a sequence of events that are controlled by several delayed internal clocks. The internal clocks are linked in such a manner that the read access time of the device is independent of the address multiplex window. $\overline{CAS}$ controls read access time; $\overline{CAS}$ must be active before or at t_{RCD} maximum to guarantee valid data out (DQ) at t_{RAC} (access time from $\overline{RAS}$ active transition). If the t_{RCD} maximum is exceeded, read access time is determined by the $\overline{CAS}$ clock active transition (t_{CAC}).

The $\overline{RAS}$ and $\overline{CAS}$ clocks must remain active for minimum times of t_{RAS} and t_{CAS}, respectively, to complete the read cycle. $\overline{W}$ must remain high throughout the cycle, and for time t_{RRH} or t_{RCH} after $\overline{RAS}$ or $\overline{CAS}$ inactive transition, respectively, to maintain the data at that bit location. Once $\overline{RAS}$ transitions to inactive, it must remain inactive for a minimum time of t_{RP} to precharge the internal device circuitry for the next active cycle. DQ is valid, but not latched, as long as the $\overline{CAS}$ clock is active. When the $\overline{CAS}$ clock transitions to inactive, the output will switch to High Z (three-state) t_{OFF} after the inactive transition.

WRITE CYCLE

The user can write to the DRAM with either an early write or a page mode early write cycle. Early write mode is discussed here, while page mode write operations are covered in a separate section.

A write cycle begins as described in **ADDRESSING THE RAM**. Write mode is enabled by the transition of $\overline{W}$ to active (V_{IL}). Early write mode is distinguished by the active transition of $\overline{W}$, with respect to $\overline{CAS}$. Minimum active time t_{RAS} and t_{CAS} and precharge time t_{RP} apply to write mode, as in the read mode.

An early write cycle is characterized by $\overline{W}$ active transition at minimum time t_{WCS} before $\overline{CAS}$ active transition. Data in (DQ) is referenced to $\overline{CAS}$ in an early write cycle. $\overline{RAS}$ and $\overline{CAS}$ clocks must stay active for t_{RWL} and t_{CWL}, respectively, after the start of the early write operation to complete the cycle.

PAGE MODE CYCLES

Page mode allows fast successive data operations at all 1024 column locations on a selected row of the module. Read access time in page mode (t_{CAC}) is typically half the regular $\overline{RAS}$ clock access time, t_{RAC}. Page mode operation consists of keeping $\overline{RAS}$ active while toggling $\overline{CAS}$ between V_{IH} and V_{IL}. The row is latched by $\overline{RAS}$ active transition, while each $\overline{CAS}$ active transition allows selection of a new column location on the row.

A page mode cycle is initiated by a normal read or write cycle, as described in prior sections. Once the timing requirements for the first cycle are met, $\overline{CAS}$ transitions to inactive for minimum t_{CP}, while $\overline{RAS}$ remains low (V_{IL}). The second $\overline{CAS}$ active transition while $\overline{RAS}$ is low initiates the first page mode cycle (t_{PC}). Either a read or write operation can be performed in a page mode cycle, subject to the same conditions as in normal operation (previously described). These operations can be intermixed in consecutive page mode cycles and performed in any order. The maximum number of consecutive page mode cycles is limited by t_{RASP}. Page mode operation is ended when $\overline{RAS}$ transitions to inactive, coincident with or following $\overline{CAS}$ inactive transition.

REFRESH CYCLES

The dynamic RAM design is based on capacitor charge storage for each bit in the array. This charge will tend to degrade with time and temperature. Each bit must be periodically **refreshed** (recharged) to maintain the correct bit state. Bits in the MCM36100 require refresh every 16 milliseconds.

This is accomplished by cycling through the 1024 row addresses in sequence within the specified refresh time. All the bits on a row are refreshed simultaneously when the row is addressed. Distributed refresh implies a row refresh every 15.6 microseconds for the MCM36100. Burst refresh, a refresh of all 1024 rows consecutively, must be performed every 16 milliseconds on the MCM36100.

A normal read or write operation to the RAM will refresh all the bits associated with the particular row decoded. Three other methods of refresh, **$\overline{RAS}$-only refresh, $\overline{CAS}$ before $\overline{RAS}$ refresh,** and **hidden refresh** are available on this device for greater system flexibility.

$\overline{CAS}$ Before $\overline{RAS}$ Refresh

This refresh cycle is initiated when $\overline{RAS}$ falls after $\overline{CAS}$ has been low (by t_{CSR}). This activates the internal refresh counter, which generates the row address to be refreshed. Externally applied addresses are ignored during the automatic refresh cycle. If the output buffer was off before the automatic refresh cycle, the output will stay in the high impedance state. If the output was enabled by $\overline{CAS}$ in the previous cycle, the data out will be maintained during the automatic refresh cycle as long as $\overline{CAS}$ is held active (hidden refresh).

Hidden Refresh

The hidden refresh method allows refresh cycles to be performed while maintaining valid data at the output pin. Hidden refresh is performed by holding $\overline{CAS}$ at V_{IL} and taking $\overline{RAS}$ high and after a specified precharge period (t_{RP}), executing a $\overline{CAS}$ before $\overline{RAS}$ refresh cycle. (See Figure 1.)

$\overline{CAS}$ BEFORE $\overline{RAS}$ REFRESH COUNTER TEST

The internal refresh counter of the device can be tested with a **$\overline{CAS}$ before $\overline{RAS}$ refresh counter test**. This refresh counter test is performed with read and write operations. During this test, the internal refresh counter generates the row address, while the external address input supplies the column address. The entire array is refreshed after 1024 test cycles, as indicated by the check data written in each row. See **$\overline{CAS}$ before $\overline{RAS}$ refresh counter test cycle** timing diagram.

The test can be performed only after a minimum of **8 $\overline{CAS}$ before $\overline{RAS}$** initialization cycles. The test procedure is as follows:

1. Write "0"s into all memory cells (normal write mode).
2. Select a column address, and read "0" out of the cell by performing **$\overline{CAS}$ before $\overline{RAS}$ refresh counter test, read cycle**. Repeat this operation 1024 times.
3. Select a column address, and write "1" into the cell by performing **$\overline{CAS}$ before $\overline{RAS}$ refresh counter test, write cycle**. Repeat this operation 1024 times.
4. Read "1"s (normal read mode), which were written at step 3.
5. Repeat steps 1 to 4 using complement data.

5

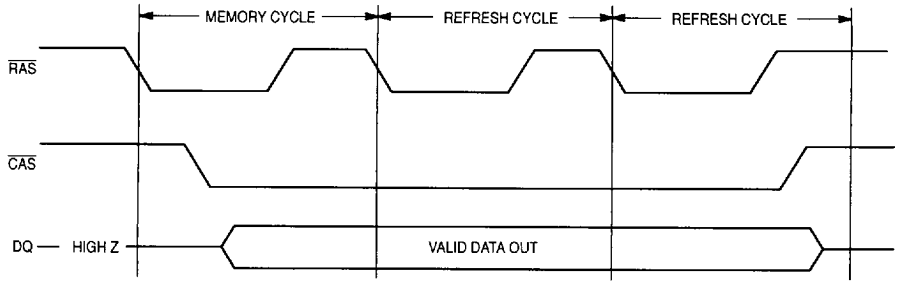

Figure 1. Hidden Refresh Cycle

ORDERING INFORMATION
(Order by Full Part Number)

MCM 36100 X XX

Motorola Memory Prefix — MCM

Part Number — 36100

Speed (60 = 60 ns, 70 = 70 ns) — XX

Package (AS = SIMM, ASG = Gold Pad SIMM, ASH = Low Profile SIMM, ASHG = Low Profile Gold Pad SIMM) — X

Full Part Numbers —

MCM36100AS60	MCM36100ASG60
MCM36100AS70	MCM36100ASG70
MCM36100ASH60	MCM36100ASHG60
MCM36100ASH70	MCM36100ASHG70